Semantik-basierte Autorenwerkzeuge für mathematische Dokumente

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Bonn, Germany 31. October 2008



1



- Support "Workflow" of preparing documents with math. content
- ... for mathematicians/scientists/software engineer for developing theories and especially proofs
- ► Why?
 - Verifiable formalisations and proofs
 - Added values through machine support
 - Take over routine checks/tasks in proofs
 - Organisation of large complex proofs (z.B. Kepler' conjecture proof by T. Hales)
 - Semantic-based search for math. concepts

▶ ...

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1 1. Simple Sets	
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Dervermon 1. (Type of Sole) Then we define the type set.	
DEFECTION 3. (Poweries \oplus) The function \oplus states are individual and a set and title whether the which belows to this set.	i inde
NOTATION 4. (Parenties \in) Let x be an independ and A a set, then we write x is independent of $A,$ is $\sin A$ or A constaining.	с A,
Deremone's. (Partice <) The function Communication labor two are end tolls whether die first set is a subset around set.	of the
Notestice 4. (Function \subset) Let A and B be sets, then we write $A \subset B$.	
$\begin{array}{l} \text{Axiout 7: } (Juginations of \subset) \\ \text{ It indications } \forall U, V, (U \subset V) \Leftrightarrow \{\forall x. \{x \in U\} \Rightarrow \{x \in V\}\} \end{array}$	
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DEFERTION 5. (Parentees U.) The function U _{mit} tasks cost index too acts and returns the scalars of both acts.	
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Axios 11. (Definition of U.)	
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Approaches

Sichere Kognitive Syste

Bottom-up Approach:

- Logic, calculus, and components build upon (proof assistants)
- Hope that eventually targeted users will use it
- Classical approach (a.o. Ωmega)



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Bottom-up Approach:

- Logic, calculus, and components build upon (proof assistants)
- Hope that eventually targeted users will use it
- Classical approach (a.o. Ωmega)

Top-down Approach:

- Start with systems that users already use
- Add functionality, e.g. those provide by proof assistants and more
- Analogy: Grammar-checker, but interactive
- Working hypotheses:
 - Authors of documents should not have to learn peculiarities of proof assistance systems in order to get their support
 - Support system should adapt to user, not vice versa
 - Author always has full control over the text (layout, formulation)

3





(Texteditors)

The Vision

Planned Functionality

Current State

Spectrum between Text, Notations and Formal Representations

Autexier: Semantik-basierte Autorenwerkzeuge für mathematische Dokumente

Bonn, October'08

Writing the Lecture Notes



Introduction to Algebra Thomas H.

1 Logic 2 Classes and Sets 3 Functions 4 Relations and Partitions ...Logic



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1 Logic

We adopt the logical foundations developed in the lecture notes [SS08, David H.] and agree on the following notational conventions:

[SS08, David H.]	This course
¬ <i>P</i>	"not <i>P</i> "
$P \wedge Q$	" <i>P</i> and <i>Q</i> "
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$\forall x.Q$	"for all x, Q"
∃ <i>z</i> . <i>Q</i>	"there exists x, Q"

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3 Functions

4 Relations and Partitions

In Gödel-Bernays form of axiomatic set theory, which we shall follow, the primitive (undefined) notions are **class**, **membership**, and **equality**. Intuitively, we consider a class to be a collection A of objects (elements) such that given any object x if it is possible to determine whether or not x is a member (or element) of A. We write $x \in A$ for "x is an element of A" and $x \notin A$ for "x is not an element of A".

[...]

The axiom of extensionality asserts that two classes with the same elements are equal (formally, $[x \in A \Leftrightarrow x \in B] \supset A = B)$.

A class \hat{A} is defined to be a **set** if and only if there exists a class B such that $A \in B$. Thus a set is a particular kind of class. A class that is not a set is called a **proper class**. Intuitively the distinction between sets and proper classes is not too clear. Roughly speaking a set is a "small" class and a proper class is exceptionnally "large". The **axiom of class formation** asserts that for any statement P(y) in the first-order predicate calculus involving a variable y, there exists a class A such that $x \in A$ if and only if x is a set and the statement P(x) is true. We denote this class A by $\{x \mid P(x)\}$.

A class A is a **subclass** of a class B (written $A \subseteq B$) provided:

for all $x \in A$, $x \in A \supset x \in B$.

By the axioms of extensionality and the properties of equality

 $A = B \Leftrightarrow A \subseteq B$ and $B \subseteq A$



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3 Functions 4 Relations and Partitions

[...]

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... Functions and Relations



Introduction to Algebra Thomas H.

1 Logic

2 Classes and Sets

3 Functions

Given classes A and B, a function (or map or mapping) f from A to B (written $f: A \rightarrow B$) assigns to each $a \in A$ exactly one element $b \in B$; b is called the value of the function at a or the image of a and is usually written f(a). A is the **domain** of the function (sometimes written Dom_i) and B is the **range** or **codomain**. Sometimes it is convenient to denote the effect of the function f on an element of A by $a \mapsto f(a)$. Two functions are **equal** if they have the same domain and range and have the same value for each element of their common domain. [...]

4 Relations and Partitions

The **axiom of pair formation** states that for any two sets [elements] *a*, *b* there is a set $P = \{a, b\}$ such that $x \in P$ if and only if x = a or x = b; if a = b then P is the **singleton** $\{a\}$. The **ordered pair** (*a*, *b*) is defined to be the set $\{\{a\}, \{a, b\}\}$; its **first component** is *a* and its **second component** is *b*. It is easy to verify that (a, b) = (a', b') if and only if a = a' and b = b'. The **Cartesian product** of classes *A* and *B* is the class

 $A \times B = \{(a, b) \mid a \in A, b \in B\}$

A subclass R of $A \times B$ is called a **relation** on $A \times B$. For example, if $f : A \to B$ is a function, the **graph** of f is the relation $R = \{(a, f(a)) \mid a \in A\}$. [...]

... Functions and Relations



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[...]

... Verification Details on Demand



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[...]

A class A is a **subclass** of a class B (written $A \subset B$) provided:

for all
$$x \in A, x \in A \Rightarrow x \in B$$
.

By the axioms of extensionality and the properties of equality $^{\ensuremath{\textit{Details}}}$

$$A = B \Leftrightarrow A \subset B$$
 and $B \subset A$

Details We first prove $A = B \Rightarrow A \subset B$ and $B \subset A$: Assume (h) A = B, then we have to prove (1) $A \subset B$ and (2) $B \subset A$: For (1), assuming $x \in A$, we conclude $x \in B$ from (h) and properties of equality. For (2), assuming $x \in B$, we conclude $x \in A$ from (h) and properties of equality. Conversely, we prove $A \subset B$ and $B \subset A \Rightarrow A = B$: By Definition of \subset we know from $A \subset B$ and $B \subset A \Rightarrow A = B$: By Definition of \subset we know from $A \subset B$ and $B \subset A \Rightarrow A = B$: By Definition of \subset we know from $A \subset B$ and $B \subset A \Rightarrow x \in B$ and $x \in B \Rightarrow x \in A$ for all x. Hence, $x \in A \Leftrightarrow x \in B$ for all x and by extensionality follows A = B.

4 Functions

1



Planned Functionality

Current State

Spectrum between Text, Notations and Formal Representations

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General

- Are all used concepts defined? Uniquely? Unambiguously?
- Have all introduced notations been followed?
- Management/Maintenance of notational information
- Use of theories from other documents (semantic citation/copy&paste)
 Specifically for proofs
- In a subproof: What are possible next steps?
- Apply automatic proof procedures (verification or subtasks)
- Automatically found (sub-)proofs
 - ▶ integrated into document: readable, e.g. for inspection, explanation
 - use introduced notations
- Is a proof complete? Is it verified?

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The Vision

Planned Functionality

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• Connect $T_E X_{macs}$ with $\Omega mega$ via Mediator $Plat\Omega$

$$T_{E}X_{macs} \longleftrightarrow Plat\Omega \longleftrightarrow \Omega mega$$

Fully annotated (manually) T_EX_{macs} document \begin{definition}[Function \$\in\$] The predicate \concept{\in}{elem \times set \rightarrow bool} takes an individual and a set and tells whether that individual belongs to this set. \end{definition}

Use Macros to indicate semantics: begin/end of theories, definitions, theorems, proofs, proof steps











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- Except for formulas.
- Provides type and proof checking, interactive proof construction
- Incremental: only changes are passed around
- Basic document change management



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Marc's Part

Next Phase: Verimathdoc II

- ► Support independent from editors (*T_EX_{macs}*, Word 2007, LAT_EX)
- Start from Document as is
- Use NL-Analysis to obtain richer structures on which support can be offered (consistency checks, reorganisation, proof support, etc.)
- Use NL-Generation to map changes and modifications back into the document.



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Formal Counterpart

Definitions



Natural Language Text

A class A is a **subclass** of a class B (written $A \subset B$) provided:

for all $x \in A, x \in A \Rightarrow x \in B$.

Formal Counterpart

```
define subclass : object*object->o;
```

Notation

notation subclass(#1, #2) <-> #1 \subset #2



The **axiom of class formation** asserts that for any statement P(y) in the first-order predicate calculus involving a variable *y*, there exists a class *A* such that $x \in A$ if and only if *x* is a set and the statement P(x) is true. We denote this class *A* by $\{x \mid P(x)\}$.

Formal Counterpart

```
axiom "class formation" ! P:object->o . ? A. class(A) /\
    ! x . membership(x,A) <=> (set(x) /\ P(x));
```

```
define setconstr:(object->o)->object;
axiom "" ! P:object->o . class(setconstr(P)) /\
        ! x . membership(x,setconstr(P)) <=> (set(x) /\ P(x));
```

Notation

notation setconstr(lam #1 . #2) <-> {#1 | #2}

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By the axioms of extensionality and the properties of equality

 $A = B \Leftrightarrow A \subset B$ and $B \subset A$

Formal Counterpart

```
conjecture "Class_Th_124" ! A,B. class(A) /\ class(B) =>
  (A = B <=> subclass(A,B) /\ subclass(B,A));
```

```
proof
fact ! A,B. class(A) /\ class(B) => (A = B <=>
    subclass(A,B) /\ subclass(B,A)) by "Extensionality", "Properties of ="
trivial
end
```



By the axioms of extensionality and the properties of equality

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Formal Counterpart

```
conjecture "Class_Th_124" ! A,B. class(A) /\ class(B) =>
     (A = B \iff subclass(A,B) / subclass(B,A));
proof
fact ! A,B. class(A) /\ class(B) => (A = B <=>
    subclass(A,B) /\ subclass(B,A)) by "Extensionality", "Properties of ="
   proof
    assume class(a), class(b);
     subgoal "1": a = b => subclass(a,b) /\ subclass(b,a);
      proof assume "hyp" : a = b;
         subgoal "1a": subclass(a,b)
           proof assume "hyp1" membership(x,a); subgoal membership(x,b);
                trivial by "hyp", "hyp1" and "=" end
         subgoal "1b": subclass(b,a)
           proof assume "hyp2" membership(x,b); subgoal membership(x,a);
                 trivial by "hyp", "hyp2" and "=" end
```

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Formal Counterpart

```
subgoal "2": subclass(a,b) /\ subclass(b,a) => a = b;
  proof assume subclass(a,b), subclass(b,a);
         fact ! x . membership(x,a) => membership(x,b) by "subclass";
         fact ! x . membership(x,b) \Rightarrow membership(x,a) by "subclass";
         fact ! x . membership(x,b) <=> membership(x,a) by "logic";
         trivial by "Extensionality"; end
end; trivial; end
```

Natural Language Text

By the axioms of extensionality and the properties of equality Details

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Formal Counterpart

```
define class:object->o;
define membership:object*object->o;
```

axiom !x,a . membership(x, a) or ~ membership(x, a);

Notation

```
notation \#1 \in \#2 \iff membership(\#1, \#2)
notation \#1 \notin \#2 \iff membership(\#1, \#2)
```